*** create a table with one column (ID) having many distinct values and one column (CODE) having very few SQL> CREATE TABLE ziggy_stuff AS SELECT mod(rownum,500000) id, mod(rownum,5) code, 'ZIGGY' name FROM dual CONNECT BY LEVEL <=1000000;

Table created.

*** However, add a row that has a very distinct CODE value. Although there are only 6 different CODE values, there's only one occurance of value 42

SQL> INSERT INTO ziggy_stuff VALUES (42, 42, 'BOWIE');

1 row created.

SQL> COMMIT;

Commit complete.

SQL> exec dbms_stats.gather_table_stats(ownname=>null, tabname=>'ZIGGY_STUFF', cascade=> true, estimate_percent=> null, method_opt=> 'FOR ALL COLUMNS SIZE 1');

PL/SQL procedure successfully completed.

*** Create a histogram on the CODE value so that the CBO knows there's very few CODEs with a value of 42 SQL> exec dbms_stats.gather_table_stats(ownname=>null, tabname=>'ZIGGY_STUFF', cascade=> true, estimate_percent=> null, method_opt=> 'FOR COLUMNS CODE SIZE 10');

PL/SQL procedure successfully completed.

*** First, create an index with the ID column being the leading column

SQL> CREATE INDEX ziggy_stuff_id_code_i ON ziggy_stuff(id, code);

Index created.

SQL> SELECT * FROM ziggy_stuff where id = 42 AND code = 42;

1 row selected.

Execution Plan

Plan hash value: 975820249

Id Operation	Name	Rows	Ī	Bytes	Ī	Cost (%CPU)	Time	Ī
0 SELECT STATEMENT 1 TABLE ACCESS BY INDEX ROWID * 2 INDEX RANGE SCAN	 ZIGGY_STUFF ZIGGY_STUFF_ID_CODE_I	1 1 1		13 13		4 (0)	00:00:01 00:00:01 00:00:01	İ

Predicate Information (identified by operation id):

2 - access("ID"=42 AND "CODE"=42)

Statistics

- 0 recursive calls db block gets consistent gets
- physical reads

- physical reads redo size bytes sent via SQL*Net to client bytes received via SQL*Net from client SQL*Net roundtrips to/from client sorts (memory) sorts (disk) rows processed

 - rows processed

*** AS expected, search on both columns and the index is used \dots

SQL> SELECT * FROM ziggy_stuff where id = 42;

3 rows selected.

Execution Plan

Plan hash value: 975820249

Id Operation	Name	Rows	Bytes	Cost (%	CPU) Time	
0 SELECT STATEMENT 1 TABLE ACCESS BY INDEX ROWID 2 INDEX RANGE SCAN	 ZIGGY_STUFF ZIGGY_STUFF_ID_CODE_I	2 2 2	26 26	6 6 3	(0) 00:00:01 (0) 00:00:01 (0) 00:00:01	1 İ

Predicate Information (identified by operation id):

2 - access("ID"=42) Statistics recursive calls recursive calls
db block gets
consistent gets
physical reads
redo size
bytes sent via SQL*Net to client
bytes received via SQL*Net from client
SQL*Net roundtrips to/from client
sorts (memory)
sorts (disk)
rows processed *** Search on only the leading column (ID) and again the index can be used effectively SQL> SELECT * FROM ziggy_stuff WHERE code = 42; 1 row selected. Execution Plan Plan hash value: 4141990364 | Id | Operation | Name | Rows | Bytes | Cost (%CPU)| Time 307 (15)| 00:00:03| 307 (15)| 00:00:03| 0 | SELECT STATEMENT | | 1 | TABLE ACCESS FULL | ZIGGY_STUFF | Predicate Information (identified by operation id): 1 - filter("CODE"=42) Statistics recursive calls
db block gets
consistent gets
physical reads
redo size
bytes sent via SQL*Net to client
bytes received via SQL*Net from client
SQL*Net roundtrips to/from client
sorts (memory)
sorts (disk)
rows processed 2605 0 rows processed *** However, search on the CODE column only and the index can not be used.
*** As the leading column is very selective, a CODE value of 42 could potentially be referenced within any of the index leaf blocks SQL> DROP INDEX ziggy_stuff_id_code_i; Index dropped.

*** Let's now re-create the index but with the columns the other way around (CODE now the leading column)

SQL> CREATE INDEX ziggy_stuff_code_id_i ON ziggy_stuff(code,id);

Index created.

SQL> SELECT * FROM ziggy_stuff WHERE id = 42 AND code = 42;

1 row selected.

Execution Plan

Plan hash value: 442388428

Id Operation	Name	Ī	Rows		Bytes	Ī	Cost	(%CPU)	Time	Ī
0 SELECT STATEMENT 1 TABLE ACCESS BY INDEX ROWID * 2 INDEX RANGE SCAN	 ZIGGY_STUFF ZIGGY_STUFF_CODE_ID_I		1 1 1	. . .	13 13		4	(0)	00:00:01 00:00:01 00:00:01	

Predicate Information (identified by operation id):

2 - access("CODE"=42 AND "ID"=42)

Statistics

- 1
- recursive calls db block gets consistent gets physical reads
- physical redus redo size bytes sent via SQL*Net to client bytes received via SQL*Net from client SQL*Net roundtrips to/from client sorts (memory) sorts (disk)

- rows processed

*** Again as expected, index is used when both columns are searched

SQL> SELECT * FROM ziggy_stuff WHERE code = 42;

1 row selected.

Execution Plan

Plan hash value: 442388428

Id Operation	Name	 Rows	Ī	Bytes	Cos	t (%CPL	1)	Time	<u> </u>
0 SELECT STATEMENT 1 TABLE ACCESS BY INDEX ROWID * 2 INDEX RANGE SCAN	 ZIGGY_STUFF ZIGGY_STUFF_CODE_ID_I	1 1 1		13 13		4 (0	Ŋί	00:00:01 00:00:01 00:00:01	

Predicate Information (identified by operation id):

2 - access("CODE"=42)

Statistics

- recursive calls
 db block gets
 consistent gets
 physical reads
 redo size
 bytes sent via SQL*Net to client
 bytes received via SQL*Net from client
 SQL*Net roundtrips to/from client
 sorts (memory)
 sorts (disk)
 rows processed

 - 000
 - rows processed

*** When searching on just the CODE column for the value 42, with the histogram in place, the CBO estimates there's only the one row and so can use the index effectively

SQL> SELECT * FROM ziggy_stuff WHERE id = 42;

3 rows selected.

Execution Plan

Plan hash value: 2304838088

Id Operation	Name	Rows	Bytes	Cost	(%CPU) Time	Ī
0 SELECT STATEMENT 1 TABLE ACCESS BY INDEX ROWID * 2 INDEX SKIP SCAN	 ZIGGY_STUFF ZIGGY_STUFF_CODE_ID_I	2 2 2	26 26		(0) 00:00:01 (0) 00:00:03 (0) 00:00:03	1

Predicate Information (identified by operation id):

2 - access("ID"=42)
 filter("ID"=42)

Statistics

- 1 recursive calls
 0 db block gets
 19 consistent gets
 10 physical reads 19 10

- physical reads redo size bytes sent via SQL*Net to client bytes received via SQL*Net from client SQL*Net roundtrips to/from client

- 0 sorts (memory) 0 sorts (disk) 3 rows processed
- *** When searching on just the ID column, the CBO knows there are only 6 distinct CODE column values
- *** The CBO can effectively probe the index in 6 different locations and retrieve all the necessary rows.
- *** At 19 consistent gets, it's not as good as the 7 consistent gets with the previous index
- *** However, it's not too bad and much better than the approx 2605 consistent gets required for a full table scan
- $\ensuremath{^{***}}$ Perhaps the second index will suffice, making the overheads associated having a second index unnecessary \dots